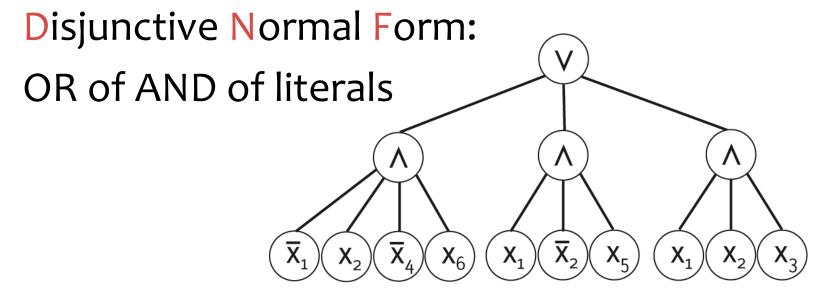
Learning Talagrand DNF Formulas

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DNF Formulas



Can also write as: $\overline{x}_1 x_2 \overline{x}_4 x_6 \lor x_1 \overline{x}_2 x_5 \lor x_1 x_2 x_3$ Size is the number of AND gates (terms).

PAC Learning DNF Formulas

A is a PAC-learner for poly(n)-size DNF if $\forall f$ in the class given uniform random examples (x,f(x)) w.h.p. outputs h s.t.

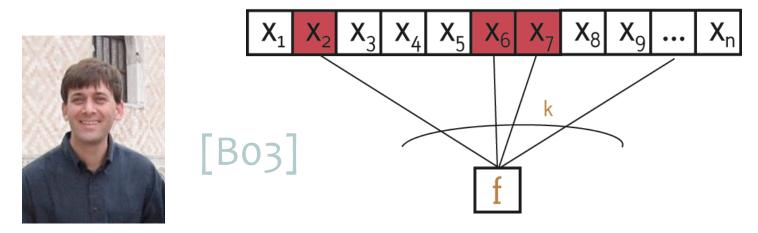
$$Pr[h(x) = f(x)] \ge 1 - \varepsilon$$
 [V84]

Best alg takes time n^{O(log n/ε)} [V90]



Juntas

Boolean funcs that depend on $\leq k$ vars.



Best alg takes time $n^{0.7k}$ [MOS03] Learning DNF \Rightarrow Learning O(log n) Juntas

Parity with Noise

$$S = \{x_1, x_5, x_8, x_9\}$$

 $\chi_{S}(x) = 1$ if odd # of vars in S are set to 1.

$$\chi_{S}(x) \oplus \eta, \eta = 1 \text{ w.p. p}$$

Best alg takes time 2^{O(n/log n)} [BKWoo]









Learning PWN, $|S|=O(\log n) \Rightarrow \text{Learning DNF}$

[FGKP06]

Statistical Queries

An SQ-oracle given g, outputs a good estimate to E[g(x,f(x))]

SQ-learners for DNF take $n^{\omega(1)}$ queries [K93]

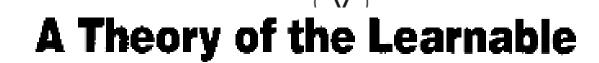


Almost all PAC-learning algs are SQ algs!

Monotone DNF

Monotone: no negations on the literals

$$X_1X_2X_4X_6 \lor X_1X_2X_5 \lor X_1X_2X_3$$

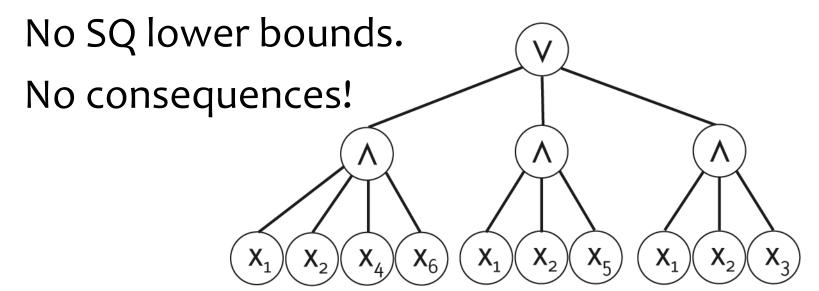


advantageous. The question as to whether monotone DNF expressions can be learned from EXAMPLES alone is open. A positive answer would be especially signifi-

No Excuses!

Monotone juntas are easy.

MDNF can't compute parity.



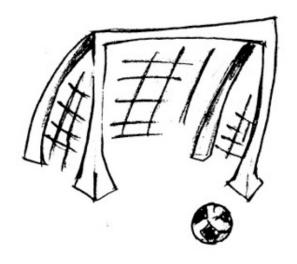
Known Results

- Poly(n)-size read-k MDNF. [HM91]
- Size-2^{√log(n)} MDNF [So1]
- Random poly(n)-size MDNF [So8,JLSWo8]
 - Pick t terms uniformly from all terms of size log(t)
 - Relies on terms not overlapping too much

Pretty pitiful.

Setting a Goal

- Read-o(1)
- Size $\Omega(n)$
- Overlapping terms



Talagrand DNF

Pick n terms from set of all terms of length log(n) defined over first log²(n) variables.

[T96]

- Size n, read-o(1).
- Know all relevant variables.
- Lots of overlap.



Talagrand DNF

Pick n terms from set of all terms of length log(n) defined over first log²(n) variables. [T96]

- f is sensitive to low noise
 Pr[f(x)≠f(y)] = Ω(1)
 y=x with each bit flipped with prob 1/log(n)
- f has high "surface area" $\Omega(\sqrt{\log(n)})$

Prizes

 PAC-learn Talagrand DNFs w.h.p. over the choice of DNF.

- PAC-learn Talagrand DNFs in the worst case.
- Prove that Talagrand DNFs require $n^{\omega(1)}$ SQ-queries [FLS10].